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Expanders: d -regular, d constant. Small $S \subseteq V$ have many neighbors $|N(S)| \geq \frac{d}{5} |S|$

Approximate $\frac{d}{n} K_n$. Pseudo-random properties. Random d -regular are expanders.

Def. A d -regular graph G is an ε -expander if $\|L_G - \frac{d}{n} L_{K_n}\| \leq \varepsilon d$

Is equivalent to $|\lambda_i - d| \leq \varepsilon d$ for all $i \geq 2$. Because $\lambda_i(\frac{d}{n} K_n) = d \quad \forall i \geq 2$.

In adjacency, is $\mu_i = d$, $|\mu_i| \leq \varepsilon d \quad \forall i \geq 2$. Ramanujan: $\varepsilon \leq \frac{2\sqrt{d-1}}{d}$

For $A, B \subseteq V$, $A \cap B = \emptyset$ $E(A, B) \triangleq \{(a, b) \in E, a \in A, b \in B\}$

Thm! For all $A, B \subseteq V$, $A \cap B = \emptyset$, $|A| = \alpha n$, $|B| = \beta n$

$|\#E(A, B) - d\alpha\beta n| \leq \varepsilon d n \sqrt{\alpha - \alpha^2} \sqrt{\beta - \beta^2}$. Close to $d\alpha\beta n$ when $\alpha\beta > \varepsilon$

Motivation: Consider choosing A and B at random, subject to $|A| = \alpha n$, $|B| = \beta n$

$\Pr[a \in A, b \in B] \approx \alpha\beta$, so $\Pr[(a, b) \in E(A, B)] \approx 2\alpha\beta$ and $\mathbb{E}[|E(A, B)|] \approx \frac{d}{2} \cdot 2\alpha\beta = d\alpha\beta n$

For K_n , $|E(A, B)| = (\alpha n)(\beta n)$, so $\omega(E_{\frac{d}{n} K_n}(A, B)) = d\alpha\beta n$.

proof of Thm 1 $\mathbb{1}_A^T L_G \mathbb{1}_B = \mathbb{1}_A^T (dI - M_G) \mathbb{1}_B = 0 - |E(A, B)| = -|E(A, B)|$

Let $H = \frac{d}{k} K_n$. $\mathbb{1}_A^T L_H \mathbb{1}_B = -\alpha \beta n$.

And, $|\mathbb{1}_A^T (L_G - L_H) \mathbb{1}_B| = \|\mathbb{1}_A^T\| \cdot \|(L_G - L_H) \mathbb{1}_B\| \leq \sqrt{\alpha n} \cdot \varepsilon d \cdot \sqrt{\beta n} = \varepsilon d \sqrt{\alpha \beta} n$

$\rightarrow (|E(A, B)| - \alpha \beta n) \leq \varepsilon d n \sqrt{\alpha \beta}$

To improve, replace $\mathbb{1}_A$ with $X_A = \mathbb{1}_A - \alpha \mathbb{1}$, $X_B = \mathbb{1}_B - \beta \mathbb{1}$, $\|X_A\| = \sqrt{\alpha - \alpha^2} n$

$X_A^T L_G X_B = \mathbb{1}_A^T L_G \mathbb{1}_B$, so $(|E(A, B)| - \alpha \beta n) \leq \varepsilon d n \sqrt{(\alpha - \alpha^2)(\beta - \beta^2)}$

For $A \subset V$, define $N(A) = \{b = \exists a \in A, (a, b) \in E\}$ ($\approx \frac{\alpha}{\varepsilon^2}$ for α small)

Tanner's Thm For $|A| = \alpha n$, $|N(A)| = \gamma n$, $\gamma \geq \frac{\alpha}{\varepsilon^2(1-\alpha) + \alpha}$ (for $\varepsilon \approx \frac{2\sqrt{\alpha\gamma}}{\alpha}$, $\approx \alpha \frac{d}{\gamma}$)

proof Let $B = V - N(A)$, so no edges between A and B . $|B| = \beta n$ $|N(A)| = \gamma n$, $\gamma = 1 - \beta$

By Thm 1: $\alpha \beta d n \leq \varepsilon d n \sqrt{\alpha \beta (1-\alpha)(1-\beta)} \Leftrightarrow \sqrt{\alpha \beta} \leq \varepsilon \sqrt{(1-\alpha)(1-\beta)} \Leftrightarrow \alpha \beta \leq \varepsilon^2 (1-\alpha)(1-\beta)$

$\Leftrightarrow \alpha(1-\beta) \leq \varepsilon^2 (1-\alpha) \gamma \Leftrightarrow \alpha \leq \varepsilon^2 (1-\alpha) \gamma + \alpha \gamma \Leftrightarrow \gamma \geq \frac{\alpha}{\varepsilon^2(1-\alpha) + \alpha}$

Ramanujan Bound (Alon-Boppana). For large n , $\lambda_2 \leq d - 2\sqrt{d-1} + \text{small}$

Want $x \perp \mathbb{1}$ so that $x^T L x = d - 2\sqrt{d-1} + \text{small}$

Idea: Let $a_i, b_i \in V$, no edges between $a \cup N(a)$ and $b \cup N(b)$

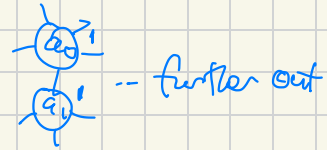


$$x(c) = \begin{cases} 1 & c=a \\ \frac{1}{\sqrt{d}} & c \in N(a) \\ -1 & c=b \\ -\frac{1}{\sqrt{d}} & c \in N(b) \end{cases}$$

$$x^T L x = 2 \left[d \left(1 - \frac{1}{\sqrt{d}}\right)^2 + d(d-1) \left(\frac{1}{\sqrt{d}}\right)^2 \right] = 2 \left[d - 2\sqrt{d} + 1 + d-1 \right] = 4(d - \sqrt{d})$$

$$x^T x = 2 \left(1 + d \frac{1}{d}\right) = 4$$

$$\frac{x^T L x}{x^T x} = d - \sqrt{d}$$



Thm If are edges (a_0, a_1) and (b_0, b_1) at dist $\geq 2k+2$, $\lambda_2 \leq d - 2\sqrt{d-1} + \frac{2\sqrt{d-1}-1}{k+1}$

proof def $A_0 = \{a_0, a_1\}$ $A_i = N(A_0) - A_0, \dots, A_k$ where $A_i = N(A_{i-1}) - \bigcup_{j < i} A_j$. Same for B.

for $a \in A_i$, $x(a) = (d-1)^{-i/2}$ $a \in B_i$, $x(a) = -\beta (d-1)^{-i/2}$. Choose β so $x \perp \mathbb{1} = 0$.

$$A\text{-denom} = \sum_{i=0}^k |A_i| (d-1)^{-i} \quad A\text{-num} = \sum_{i=0}^{k-1} |A_i| (d-1) \left((d-1)^{-i/2} \left(1 - \frac{1}{\sqrt{d-1}}\right) \right)^2 + |A_k| (d-1) (d-1)^{-k}$$

$$(d-1) \left(1 - \frac{1}{\sqrt{d-1}}\right)^2 = d - 2\sqrt{d-1}$$

$$= \sum_{i=0}^{k-1} |A_i| \frac{d - 2\sqrt{d-1}}{(d-1)^i} + \frac{|A_k| (d-1)}{(d-1)^k} = \sum_{i=0}^k |A_i| \frac{d - 2\sqrt{d-1}}{(d-1)^i} + |A_k| \frac{2\sqrt{d-1} - 1}{(d-1)^k}$$

$$\sum_{i=0}^k |A_i| \frac{d - 2\sqrt{d-1}}{(d-1)^i} \leq (d - 2\sqrt{d-1}) \text{ A-denom and } |A_k| = (d-1)^{k-i} |A_i| \rightarrow |A_k| \leq \frac{1}{k+1} \sum_{i=0}^k |A_i| (d-1)^{k-i}$$

$$\rightarrow |A_k| \frac{2\sqrt{d-1}-1}{(d-1)^k} \leq \frac{2\sqrt{d-1}-1}{k+1} \sum_{i=0}^k |A_i| (d-1)^{-i} \leq \frac{2\sqrt{d-1}-1}{k+1} \text{ (A-denom)}$$

So, $\frac{x^T L x}{x^T x} = \frac{\text{A-numer} + \text{B-numer}}{\text{A-denom} + \text{B-denom}} \leq \max \left(\frac{\text{A-numer}}{\text{A-denom}}, \frac{\text{B-numer}}{\text{B-denom}} \right) \leq d - 2\sqrt{d-1} + \frac{2\sqrt{d-1}-1}{k+1}$

