H is an 2-approx of G if

$$(1-2) G \leq H \leq (1+2)G$$
weighted expender is 2-approx of complete graph
Today every weighted graph G has an E-approx H
with \in 4 m lum edges
Implications. As $1s^{T}Le 1s = w(\partial(s))$
toandary of every set has similar weight.
 $1s^{T}_{1sa}Le 1sa = d(a)$.
So weighted degrees are similar.
 $=s$ if have fewer edges they weet have
higher weight
Then (next weet) $L_{H} \leq L_{G} \leq > L_{G}^{4} \leq L_{H}^{4}$
so inverses similar
Will have edges of $H \leq edges$ of G.
 $=s$ cat edges weist be in H O O
 E_{X} . $a = \frac{1-1}{2}b^{T}_{1}k = k paths legth left$

If
$$l > k$$
, must include edge (a16)
Consider function
 $a = \frac{1}{2} = \frac{2}{k!} + \frac{k!}{k!}$
Sum is $\frac{k^2}{a!b} + \frac{(k!)l \cdot 1^2}{peths}$
if $k > l_1$ most of energy is on this edge.
Random sampling: set Paib for each edge allo
include with prob Prich,
if do so, make weight $\frac{w_{a,b}}{P_{a,b}}$
 $L_G = \sum_{(a,b) \in E} w_{a,b} h_{a,b}$
 $E = L_H = \sum_{(a,b) \in E} P_{a,b} \frac{w_{a,b}}{P_{a,b}} h_{a,b} = L_G$

Need to bound deviations

Standard: Chernoff Boxends
For random variables
$$\times \dots \times n \in [0, \mathbb{R}]$$
 $\mu = \mathbb{E} \Xi \times i$
 $\mathbb{P} \left[\Xi \times i \ge (1 + \varepsilon) \mu \right] \le e^{-\frac{\varepsilon^2 \mu}{3\mathbb{R}}}$
 $\mathbb{P} \left[\Xi \times i \le (1 - \varepsilon) \mu \right] \le e^{-\frac{\varepsilon^2 \mu}{2\mathbb{R}}}$
 $Matrix \quad Chernoff$
For random PSD matrices $\times \dots, \times m$, $\|1 \times i\| \in \mathbb{R}$
 $M = \mathbb{E} \Xi \times i$
 $P = \left[M_{max} \left(\Xi \times i \right) \ge (1 + \varepsilon) M_{max}(M) \right] \le n e^{-\frac{\varepsilon^2}{2\mathbb{R}}}$
 $\mathbb{P} \left[M_{max} \left(\Xi \times i \right) \ge (1 + \varepsilon) M_{max}(M) \right] \le n e^{-\frac{\varepsilon^2}{2\mathbb{R}}}$

$$\frac{T \operatorname{ransform}}{For \operatorname{invertible} C, \quad A \neq B \leq > CAC^{T} \notin CBC^{T}}$$
So, $L_{H} \notin (I + \varepsilon) L_{G} \neq = L_{G} L_{H} L_{G}^{+1/2} \notin L_{G} L_{G} L_{G}^{+1/2}$

$$= TI = \frac{1}{m} L_{K_{m}}$$

$$T \text{ is projection, identity orthogonal to 1}.$$

$$E L_{H} = L_{G} \implies E L_{G} L_{H} L_{G}^{-2} = TI$$
Apply Matrix Character to this, ignoring span (1)

So,
$$M_{\text{max}}(M) = 1$$
, and onth 1, $M_2(M) = 1$

$$X_{a,b} = \begin{cases} \frac{\omega_{a,b}}{R_{a,b}} & \frac{4}{L_{a}} & L_{a,b} & L_{a} \\ 0 & 0.00 \end{cases}$$

$$Choose P_{a,b} = \frac{1}{R} \quad \omega_{a,b} \left[\left| L_{a} & L_{a,b} & L_{a} \\ L_{a} & L_{a} & L_{a} \\ So_{1} & \left[\left| X_{a,b} \right| \right] \\ = R \\ Best bound when make this center m \\ Try R = \frac{\epsilon^{2}}{3.5 \ln n} \\ Then error prob is = 2ne^{\frac{-\epsilon^{2}}{3R}} = 2ne \\ = 2n^{-\frac{16}{3R}} = 0 \end{cases}$$

Pail = I · Waib · Reff (a, b)

Thun
$$\sum_{(n,k) \in E} u_{n,k} \operatorname{Reff}(a_{n}b) = n-1$$

$$= \sum_{a \sim b} u_{n,k} (\overline{a}_{a} - \overline{b}_{b})^{T} L_{6}^{+} (\overline{b}_{a} - \overline{b}_{b})$$

$$= \sum_{a \sim b} u_{n,k} \operatorname{Tr} \left[(\overline{b}_{a} - \overline{b}_{b})^{T} L_{6}^{+} (\overline{b}_{a} - \overline{b}_{b})^{T} L_{6}^{+} \right]$$

$$= \sum_{a \sim b} u_{n,b} \operatorname{Tr} \left[(\overline{b}_{a} - \overline{b}_{b}) (\overline{b}_{a} - \overline{b}_{b})^{T} L_{6}^{+} \right]$$

$$= \operatorname{Tr} \left[\left(\sum_{a \sim b} u_{n,k} (\overline{b}_{a} - \overline{b}_{b}) (\overline{b}_{a} - \overline{b}_{b})^{T} \right) L_{6}^{+} \right]$$

$$= \operatorname{Tr} \left[L_{6} L_{6}^{+} \right] = n-1$$

$$\Longrightarrow \mathbb{E} \# \operatorname{edges} = \frac{n-1}{\mathbb{P}} \leq \frac{3.5 \operatorname{n} \ln n}{\varepsilon^{2}}$$
Chernoff says unlikely to be more than $\frac{4 \operatorname{ulms}}{\varepsilon^{2}} \operatorname{edses}.$

What if Parts > 1? Set Parts = Min (1, ± Wards Reff(arb)) include edge (arb) with prots 1, redo analysis on rest of graph without this edge. or, subdivide the edge into TParts Parallel edges by dividing weight equally. Now, all Parts = 1, and little increase in E # edges.

Note can quickly estimate Reff!

New technique: short cide decomposition.