

A d -regular graph G is an ε -expander

if $|\mu_2| \leq \varepsilon d$ and $|\mu_n| \geq -\varepsilon d$

Equivalent to $(\lambda_i - d) \leq \varepsilon d$ for $i=2$, Laplacian

And $\|L_G - \frac{d}{n}K_n\| \leq \varepsilon d$

where $\|A\| = \max_x \frac{\|Ax\|}{\|x\|}$

for sym A , $\|A\| = \max |\text{eigval}(A)|$

recall for $i \geq 2$, $\lambda_i(K_n) = n$, so $\lambda_i\left(\frac{d}{n}K_n\right) = d$

How small ε ?

Ramanujan: $\varepsilon \leq \frac{2\sqrt{d-1}}{d}$

Today: $\varepsilon \leq \frac{2\sqrt{d-1}}{d} - c$, $c > 0$, impossible for large n , d fixed.

For $A, B \subseteq V$, $A \cap B = \emptyset$

$E(A, B) \triangleq \{(a, b) \in E, a \in A, b \in B\}$

If G is d -regular, pick A and B at random, sizes $\alpha n, \beta n$

$$\Pr[a \in A, b \in B] = \alpha \beta$$

$$\Pr[(a, b) \in E(A, B)] = 2\alpha \beta$$

$$|E(E(A, B))| = \frac{d n}{2} \cdot 2\alpha \beta = d\alpha \beta n$$

For K_n , $|A|=x_n$, $|B|=y_n$

$$E(A, B) = (\alpha \eta)(\beta n) = \alpha \beta n^2, \text{ but } d = n -$$

Then If A, B s.t. $A \cap B = \emptyset$,

$$|E(A, B) - dx_n| \leq \varepsilon d n \sqrt{\alpha - \alpha^2} \sqrt{\beta - \beta^2}$$

is small error when $\alpha \beta > \varepsilon$

Proof $1_A^T L_G 1_B = -E(A, B)$

$$\text{let } t = \frac{d}{n} K_n \quad 1_A^T L_H 1_B = -\frac{d}{n} |A| \cdot |B| = d \alpha \beta n$$

$$|1_A^T (L_G - L_H) 1_B| = |E(A, B) - d \alpha \beta n|$$

$$\leq \|1_A\| \cdot \| (L_G - L_H) 1_B \| \quad \text{is Cauchy-Schwarz}$$

$$\leq \|1_A\| \cdot \|L_G - L_H\| \cdot \|1_B\| \quad \text{by norm}$$

$$\leq \sqrt{\alpha} \sqrt{\beta n} \varepsilon d = d \sqrt{\alpha} \sqrt{\beta} \varepsilon n$$

to improve, let $x_A = 1_A - \alpha \mathbf{1}$ $x_B = 1_B - \beta \mathbf{1}$,

$$\text{now, } 1_A^T (L_G - L_H) 1_B = x_A^T (L_G - L_H) x_B,$$

$$\|x_A\| = \sqrt{n(\alpha - \alpha^2)} \quad \|x_B\| = \sqrt{n(\beta - \beta^2)}$$

gives the result.

For $A \subset V$, $N(A) = \{b : \exists a \in A, (a, b) \in E\}$

Tanner's Theorem

For $|A| = \alpha n$, $|N(A)| = \gamma n$, $\gamma \geq \frac{\alpha}{\varepsilon^2(1-\alpha) + \alpha}$

For α small, think of α as $\frac{\alpha}{\varepsilon^2}$ for $\varepsilon \approx \frac{\sqrt{d-1}}{\alpha}$

$$\text{is like } \alpha \frac{d^2}{4(d-1)} \approx \alpha \frac{d}{4}$$

Proof: $T = U - N(A)$ so no edges between T and A

$$\beta_n = |T|. \quad \beta = 1 - \gamma$$

$$\alpha \beta d n \leq \varepsilon d n \sqrt{\alpha(1-\alpha)\beta(1-\beta)}$$

$$\sqrt{\alpha\beta} \leq \varepsilon \sqrt{(1-\alpha)(1-\beta)}$$

$$\alpha\beta \leq \varepsilon^2 (1-\alpha)(1-\beta)$$

$$\frac{\beta}{1-\beta} \leq \frac{\varepsilon^2(1-\alpha)}{\alpha}$$

$$\frac{1-\gamma}{\gamma} \leq \frac{\varepsilon^2(1-\alpha)}{\alpha}$$

$$\frac{1}{\gamma} \leq \frac{\varepsilon^2(1-\alpha)}{\alpha} + 1$$

$$\gamma \geq \frac{\alpha}{\varepsilon^2(1-\alpha) + \alpha}$$

How small can ε be?

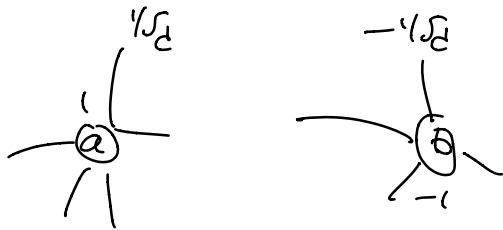
For $|A|=1$, $d=\frac{1}{n}$, $|N(A)|=d$, $\alpha=\frac{d}{n}$, $\varepsilon \geq \sqrt{d}$

Will show for large graphs is an $\times 1$ s.t.

$$\frac{x^T L x}{x^T x} \leq d - 2\sqrt{d-1} + \text{small}$$

Let $a, b \in V$ $N(a) \cap N(b) = \emptyset$
and no edges between $N(a)$ and $N(b)$

$$x(c) = \begin{cases} 1 & \text{if } c=a \\ -1 & \text{if } c=b \\ \sqrt{d} & c \in N(a) \\ -\sqrt{d} & c \in N(b) \end{cases}$$



$$x \perp \mathbb{1}, \text{ so } \lambda_2 \leq \frac{x^T L x}{x^T x}$$

$$= \frac{d\left(1 - \frac{1}{\sqrt{d}}\right)^2 + d(d-1)\left(\frac{1}{\sqrt{d}}\right)^2 + \text{some}}{\text{some} + d\left(\frac{1}{\sqrt{d}}\right)^2 + \text{some}}$$

$$= \frac{d - 2\sqrt{d} + 1 + d - 1}{2} = d - \sqrt{d}$$

We already knew that!

Improve by basing α_S on far apart edges

Assume we edges $(a_0, a_1), (b_0, b_1)$ at distance $\geq 2k+2$

If d fixed, $n \rightarrow \infty$, this happens.

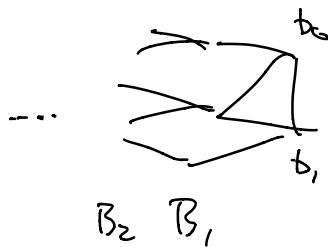
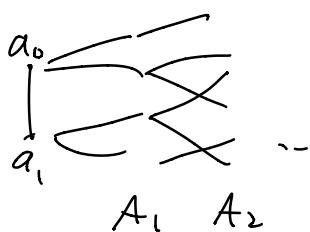
$$\lambda_2 \leq d - 2\sqrt{d-1} + \frac{2\sqrt{d-1}-1}{k+1}$$

$$A_0 = \{a_0, a_1\} \quad A_i = \{a : \text{dist}(a, A_0) = i\}$$

$$B_0 = \{b_0, b_1\} \quad B_i = \{b : \text{dist}(b, B_0) = i\}$$

A_0, \dots, A_k disjoint, B_0, \dots, B_k disjoint

no edges from A_i to B_j , any $i, j \leq k$



$$x(A_0) = 1$$

$$x(A_i) = (d-1)^{-i/2}$$

$$x(B_0) = -\beta$$

$$x(B_i) = -\beta (d-1)^{i/2}$$

choose β s.t. $\mathbf{1}^T x = 0$

$$\frac{x^T L x}{x^T x} = \frac{A\text{-edges} + B\text{-edges}}{A\text{-verts} + B\text{-verts}} \leq \max \left(\frac{A\text{-edges}}{A\text{-verts}}, \frac{B\text{-edges}}{B\text{-verts}} \right)$$

$$A\text{-verts} = \sum_{i=0}^k |A_i| (d-i)^{-\frac{i}{2}}$$

$$A\text{-edges} \leq \sum_{i=0}^{k-1} |A_i| (d-i) \frac{\left(1 - \frac{1}{\sqrt{d-i}}\right)^2}{\left((d-i)^{-\frac{i}{2}}\right)^2} + |A_k| (d-i) \cdot (d-i) (d-i)^{-k}$$

$$= \sum_{i=0}^{k-1} \frac{|A_i|}{(d-i)^i} (d-2\sqrt{d-1}) \quad \begin{aligned} & (d-i) \left(1 - \frac{1}{\sqrt{d-i}}\right)^2 \\ &= (\sqrt{d-i} - 1)^2 = d-1 - 2\sqrt{d-1} + 1 \\ &= d-2\sqrt{d-1} \end{aligned}$$

$$+ \frac{|A_k|}{(d-i)^k} (d-i)$$

$$= \underbrace{\sum_{i=0}^k \frac{|A_i|}{(d-i)^i} (d-2\sqrt{d-1})}_{\leq (d-2\sqrt{d-1}) A\text{-verts}} + \underbrace{\frac{|A_k|}{(d-i)^k} (2\sqrt{d-i} - 1)}_d$$

$$|A_k| \leq (d-i)^{k-i} |A_i|$$

$$\Rightarrow |A_k| \leq \frac{1}{k+1} \sum_{i=0}^k (d-i)^{k-i} |A_i|$$

$$\frac{|A_k|}{(d-i)^k} (2\sqrt{d-i} - 1) \leq \frac{1}{k+1} (2\sqrt{d-i} - 1) \cdot \underbrace{\sum_{i=0}^k \frac{|A_i|}{(d-i)^i}}_{\leq A\text{-verts}}$$