

Rayleigh's Monotonicity Theorem:

increasing resistances does not decrease effective resistance.

Do with effective spring constants = $\frac{l}{R_{\text{eff}}}$

Let $C_{\text{eff}}(s,t)$ = effective spring constant of $G = (V, E, w)$
between s and t

$$= \min_{\substack{x(s)=0 \\ x(t)=1}} x^T L x$$

Let $\hat{G} = (V, E, \hat{w})$ satisfy $\hat{w}_{a,b} \leq w_{a,b}$ for all $(a,b) \in E$.

Let $\hat{C}_{\text{eff}}(s,t)$ = version for \hat{G} . \hat{L} Lap of \hat{G} .

Thm $\forall s, t \quad \hat{C}_{\text{eff}}(s,t) \leq C_{\text{eff}}(s,t)$

proof

$$C_{\text{eff}}(s,t) = \min_{\substack{x(s)=0 \\ x(t)=1}} x^T L x.$$

Let x be vector on which minimum is achieved.

$$\text{Then } x^T L x \geq x^T \hat{L} x \geq \min_{\substack{\hat{x} \text{ s.t.} \\ \hat{x}(s)=0 \\ \hat{x}(t)=1}} \hat{x}^T \hat{L} \hat{x} = \hat{C}_{\text{eff}}(s,t).$$

Demo?

For an edge (a,b) , $\text{Reff}(a,b)$ is a measure of its importance. Will later sample edges with prob proportional to $\text{Reff}(a,b)$.

Let's see how to compute and store estimator efficiently.

$$\begin{aligned}\text{Recall } \text{Reff}(a,b) &= (\delta_a - \delta_b)^T L^+ (\delta_a - \delta_b) \\ &= \|L^{+1/2} (\delta_a - \delta_b)\|^2 \\ &= \|L^{+1/2} \delta_a - L^{+1/2} \delta_b\|^2 \\ &= \text{dist}(L^{+1/2} \delta_a, L^{+1/2} \delta_b)^2\end{aligned}$$

Is square of a Euclidean distance.

Johnson-Lindenstrauss:

A Euclidean distance on n vectors x_1, \dots, x_n is well-approximated by a distance in $O(\epsilon^{-2})$ dimension.

Will use Gaussian random variables.

Thm Let $x_1, \dots, x_n \in \mathbb{R}^k$, $\epsilon > 0$, $\delta > 0$. $d = \frac{8 \ln(n^2/\delta)}{\epsilon^2}$.

Let R be random d -by- k matrix of independent $N(0, 1/d)$ variables.

Then with prob $\geq 1 - \delta$, $\forall a \neq b$

$$(1 - \epsilon) \text{dist}(x_a, x_b) \leq \text{dist}(Rx_a, Rx_b) \leq (1 + \epsilon) \text{dist}(x_a, x_b)$$

So, instead of storing $x_a = L^{+1/2} \delta_a \in \mathbb{R}^n$

store $y_a = R L^{+1/2} \delta_a \in \mathbb{R}^d$

Computing $(\delta_a - \delta_b)^T L^+ (\delta_a - \delta_b)$ requires solving a system in L .

Note can do this in time essentially $O(m \sqrt{gn})$.

Will see how to approx all $\text{Reff}(a,b)$ using only $O(gn)$ such solves.

Recall $L = U W U^T$ where $U \in \mathbb{R}^{n \times m}$ signed edge-vertex
 $W \in \mathbb{R}^{m \times m}$ diagonal of weights

$$\begin{aligned} \text{We have } L^+ &= L^+ L L^+ = L^+ U W U^T L^+ \\ &= L^+ U W^{1/2} \cdot W^{1/2} U^T L^+ \end{aligned}$$

$$\begin{aligned} \text{so, } (\delta_a - \delta_b)^T L^+ (\delta_a - \delta_b) &= \left\| W^{1/2} U^T L^+ (\delta_a - \delta_b) \right\|^2 \\ &= \text{dist}(M \delta_a, M \delta_b)^2 \end{aligned}$$

$$\text{for } M = W^{1/2} U^T L^+$$

Choose $R \in \mathbb{R}^{d \times m} \sim N(0, \frac{1}{\sqrt{d}})$, and set

$$y_a = R M \delta_a \in \mathbb{R}^d$$

By JL, probably $\text{dist}(y_a, y_b)^2 \in (1 \pm \epsilon) \text{Reff}(a,b)$, $\forall a,b$.

To compute, multiply each of d rows of R
by $W^{1/2}$, U^T and then L^T
 $O(m)$ entries can do quickly

Gaussian Random Variables

$N(0,1)$ has density $p(x) = \frac{1}{\sqrt{2\pi}} e^{-x^2/2}$

If $X = \sum_{i=1}^n \tau_i$ $\tau_i \in \pm 1$ prob $1/2$

$$\Pr \left[\frac{X}{\sqrt{n}} \in [a,b] \right] \rightarrow \int_a^b p(x) dx$$

$N(0, \sigma^2)$ has density $\frac{1}{\sqrt{2\pi} \sigma} e^{-x^2/2\sigma^2}$

multiply an $N(0,1)$ by σ .

Claim If τ_1, \dots, τ_n are indep $\tau_i \leftarrow N(0, \sigma_i^2)$,

then $\sum \tau_i$ is $N(0, \sum \sigma_i^2)$

Variances of sums of indep vars add.

This says is Gaussian.

For $x \in \mathbb{R}^d$ and τ a vector of indep $N(0,1)$ vars,

$x^T \tau$ is $N(0, \|x\|^2)$

because $x^T r = \sum_i x(i) r(i)$
 x is $\mathcal{N}(0, x(i)^2)$

If r is vector of indep $\mathcal{N}(0, \sigma^2)$ vars,
 $x^T r$ is $\mathcal{N}(0, \sigma^2 \|x\|^2)$.

let r be a vector of d indep $\mathcal{N}(0, 1)$ vars.
 $\|r\|^2$ is called a χ^2 random variable.

Thm For $\varepsilon < 1$,

$$\Pr[|\|r\|^2 - d| > \varepsilon d] \leq 2 \exp(-\varepsilon^2 d / 8)$$

Now, to prove JL.

Thm let $x_1, \dots, x_n \in \mathbb{R}^k$, $\varepsilon > 0$, $\delta > 0$. $d = \frac{8 \ln(n^2 / \delta)}{\varepsilon^2}$.

let R be random d -by- k matrix of
independent $\mathcal{N}(0, 1/d)$ variables.

Then with prob $\geq 1 - \delta$, $\forall a \neq b$

$$(1 - \varepsilon) \text{dist}(x_a, x_b) \leq \text{dist}(Rx_a, Rx_b) \leq (1 + \varepsilon) \text{dist}(x_a, x_b) \quad (*)$$

proof for $a \neq b$, (*) fails when

$$\left| \|R x_a - R x_b\|^2 - \|x_a - x_b\|^2 \right| > \varepsilon \|x_a - x_b\|^2$$

$$\|R x_a - R x_b\|^2 = \|R(x_a - x_b)\|^2.$$

Each entry of $R(x_a - x_b)$ is $N\left(0, \frac{\|x_a - x_b\|^2}{d}\right)$

Translating theorem says

$$\mathbb{P}\left[\left|\|R(x_a - x_b)\|^2 - \|x_a - x_b\|^2\right| > \varepsilon \|x_a - x_b\|^2\right] \leq 2 \exp\left(-\frac{d \varepsilon^2}{8}\right)$$

$$= 2 \exp\left(-\ln(n^2/\delta)\right) = \frac{2\delta}{n^2}$$

So, prob $\exists a, b$ that violate (*)

$$\leq \binom{n}{2} \frac{2\delta}{n^2} < \delta.$$