Effective resistance
$$V = IR$$
 $R = \frac{V}{I}$

$$= \frac{\text{udtage difference}}{\text{carrent flow.}}$$

Consider lext =
$$\delta_a - \delta_b$$
 flows 1 from a to b in graph

 $N = L^{\dagger} i_{ext} = L^{\dagger} (\delta_a - \delta_b)$
 $difference = N(a) - N(b) = (\delta_a - \delta_b)^T N$
 $Reft (a,b) = (\delta_a - \delta_b)^T L^{\dagger} (\delta_a - \delta_b)$

Will see later that is a distance. For now,

let
$$L^{th} = (L^t)^{'l_2}$$

Every psol matrix A has $A^{'l_2}$ s.t. $A^{'l_2} = A$

if $A = \sum_i \lambda_i \, \psi_i \, \psi_i^T$, $A^{'l_2} = \sum_i \lambda_i \, \psi_i \, \psi_i^T$

Reff(a1b) =
$$(\int_{a}^{a} - \int_{b}^{a})^{T} L^{+12} \cdot L^{+12} (\int_{a}^{a} - \int_{b}^{b})^{T}$$

= $\|L^{+12} (\int_{a}^{a} - \int_{b}^{+12})\|_{2}^{2}$
= $\|L^{+12} \int_{a}^{a} - L^{+12} \int_{b}^{a}\|_{2}^{2}$

By energy minimization.

Recall in spring of constant ω , $\mathcal{E} = \frac{1}{2}\omega l^2$ when stretch to length l.

So, let's stretch to less 1,
measure minimum energy, and double to ret
effective spring constant

 $\mathcal{L}(x) = \frac{1}{2} \sum_{a \neq b} W_{a(b)} (x(a) - x(b))^2$

lets for ×(s)=1, ×(t)=0

And want harmonic else whose.

Set $\gamma = \frac{L^{+}(\delta_{s} - \delta_{t})}{\text{Reff(s,t)}}$

Gues 4(s)-1(t)= (5s-5t) (1s-5t) =1

Could shift by x= y- 1y/t/

Now x(t)=0, x(s)=1

Hermonic on U- {s,+? => minumes every

$$\mathcal{E}(x) = \frac{1}{2} x^{T} | x = \frac{1}{2} x^{T} | x$$

$$= \frac{1}{2} \frac{1}{Reff(r_{t}t)^{2}} (\delta_{s} - \delta_{t})^{T} | t^{T} | t^{$$

So, effective spring constant = PROTECSIT)

Classic examples

Path with n vertices, edges of resistance Ti...Thy

Reff(Iin) = Titi-+Th-1

proof set $N(a) = T_1 + \cdots + T_{a-1}$ V(i) = 0, $V(n) = T_1 + \cdots + T_{a-1}$ convert over edge $(a \cdot i, a)$ is $\underbrace{V(a \cdot i) - V(o)}_{T_{a-1}} = \underbrace{-T_{a-1}}_{T_{a-1}} = 1$

So, corresponds to a flow of value I from u to 1.

Parallel edges

Clain Reff (s,t) = 1/t, +··+ 4cn

Proof
$$R = \frac{U}{T}$$
 sof $v(s) = (v(t) = 0)$

flow on edge i is $(u(s) - v(t)) = \frac{1}{T_i}$

So, total flow =
$$\frac{7}{2}$$
 $\frac{1}{7}$

= $\frac{1}{7}$
 $\frac{1}{7}$

If view formula as $\hat{i} = (u(\hat{s}) - v(t)) \omega_{s,t}$ $\omega_{\hat{i}} = /r_{\hat{i}}$, then add weights of parallel edges.

Equivalent networks. Given B, went matrix LB St. iB=LBN(B) when N hormoniz on S=V-B.

To do it slowly, first consider $B = \{2,...,n\}$ $S = \{i\}$ Let $N = \{a: a \sim 1\}$.

Want to compute LN given N(B) and N(i) = III) Zwing N(o)

Substitute to U(i) in

Vert (a) = J(a) N(a) - Z Wall N(b) when 149

no charge if at 1

For and becomes

$$\frac{d(a) \, \mathcal{N}(a)}{b \sim q} - \sum_{b \neq 1} \omega_{a_{1}} \, \mathcal{N}(b) - \frac{\omega_{a_{1}1}}{d^{(1)}} \sum_{c \sim 1} \omega_{i,c} \, \mathcal{N}(c)$$

$$= \mathcal{N}(\hat{\alpha}) \left[d(\hat{\alpha}) - \frac{\omega_{i,4}^{2}}{du} \right] - \sum_{\substack{C \sim i \\ C \neq a}} \frac{\omega_{a_{i1}} \omega_{i_{1}c}}{du} \mathcal{N}(c) - \sum_{\substack{b \sim a \\ b \neq i}} \omega_{a_{ib}} \mathcal{N}(b)$$

Claim is result of elimination on roulcol 1 and, is a Loplacian equation on B

We removed node I and attacked edges, and pat back a clique on nbs, where for a,c~1 have edge of wt wo, we,1

proof is loplairen

(i is symmetr: some chare to iest(c) in a

2. of diagonal torms nosative

3. sum of coefficients is 0, as

$$W_{l,a} - \frac{W_{l,a}^2}{dll} - \sum_{\substack{C \sim l \\ C \neq q}} \frac{w_{l,c} w_{l,q}}{dll} = W_{l,a} - \frac{w_{l,q}}{dll} \sum_{C \sim l} w_{l,c} = 0$$

$$V(B)^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \end{pmatrix}^{T} \downarrow_{A} \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega_{i,\alpha} v(\alpha) \\ v(B) \end{pmatrix}^{T} \downarrow_{B} V(B) = \begin{pmatrix} \frac{1}{4} & \sum_{\alpha \in A} \omega$$

idea:
$$\frac{1}{dl}\sum_{\alpha \sim l} \omega_{l,\alpha} \nu(\alpha) = -\frac{L(l,\beta) \nu(\beta)}{L(l,l)}$$

substituting this in yields

$$V(B)^{T}$$
 $\left[L(B,B) - \frac{L(B,I)L(IIB)}{LB}\right]V(B)$

To check Laplacian, note: only decrease entires,

Or,
$$11_B^T L(B_1B) 11_B = d(i)$$
 and $L(I_1B) 11_B = d(i)$
 $L(I_1I) = d(i)$.

Is what get by GE on rov(rol1.

Eliminating many vertices at once:

Does not depend on order!

To elim entres in row aeB and cols in S using rows in S, mult by coeffs C so that $L(a_1S) - CL(S_1S) = O$ So, $C = L(a_1S) L(S_1S)^{-1}$

giving Lola, = L(a,:) - L(a,s) L(s,s) [L(a,:)

restricting to rows and cols in B we set

LB(B,B) = $L(B,B) - L(B,S)L(S,S)^{-1}L(S,B)$ is Schur Complement with respect to S

or onto B

To show equiv of harmons on S:

harmons => L(s,s)v(s) + L(s,B)v(B) = 0 $v(s) = -L(s,s)^{-1}L(s,B)v(B)$ iext(B) = L(B,S)v(s) + L(B,B)v(B) $= [L(B,B) - L(B,S)L(s,S)^{-1}L(s,B)]v(B)$

If B={sit} get down to one edge, whose weight is Treff(sit)

In particula, for s, t & B,

 $(\delta_s - \delta_t)^T L_B^{\dagger} (\delta_s - \delta_t) = (\delta_s - \delta_t)^T L_B^{\dagger} (\delta_s - \delta_t)$

Is how GE works to solve iext=Lv

Order ventices (,.., n

Construct Lianing for each a.

Given iest (a) and N(a+i), , V(n)

solve for N(a).

If iest (a) = 0, N(a) = 10 \ bra

O(w need to account for ivent (a)

Reff as distance:

assert Yaibic Reff(aib) + Reft(bic) > Roff(aic)

Only need to prove for 3-node graphs.