

$$\underbrace{\quad}_{x_1} + \underbrace{\quad}_{x_2} + \underbrace{\quad}_{x_3} + \underbrace{\quad}_{x_4} + \underbrace{\quad}_{x_5} = 3$$

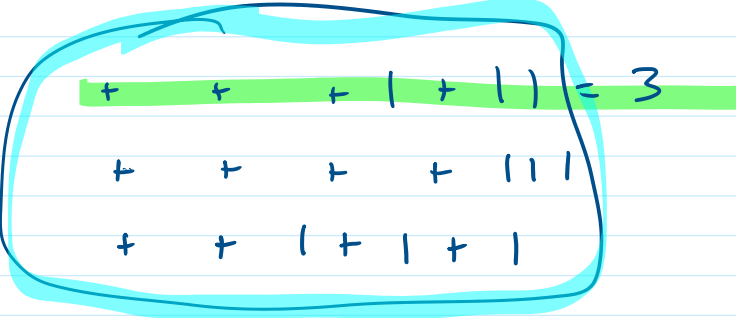
cherry # grape # lime # orange # lemon

$$0 \leq x_i \leq 3$$

$$0 + 0 + 0 + 1 + 2 = 3$$

$$0 + 0 + 0 + 0 + 3 = 3$$

$$0 + 0 + 1 + 1 + 1 = 3$$



strings of
7 l's and '+'s
w/ 4 r's and 3 l's

total # of such strings = total ways of choosing $x_1, \dots, x_5 = \binom{7}{3} = \binom{7}{4}$
 $\binom{n}{k} = \binom{n}{n-k}$

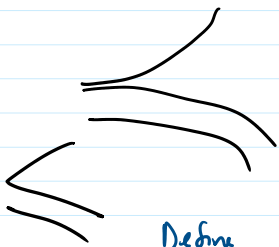
Define R on $\mathcal{P}(\mathbb{N})$ by $A R B$ iff $A \subseteq B$

	reflexive	YES	$\forall A \subseteq \mathbb{N}, A \subseteq A$
$A \subseteq B$ and $B \subseteq A$ only when $A=B$	symmetric	NO	<u>antisymmetric</u> $\forall a, b \in A \rightarrow a R b \wedge a \neq b \rightarrow b \not R a$
	transitive	YES	$\forall A, B, C \subseteq \mathbb{N}, A \subseteq B \wedge B \subseteq C \rightarrow A \subseteq C$

$a \leq b \wedge a \neq b \rightarrow \neg(b \leq a)$

DEF: A partial order \leq on a set S is a relation on S that is reflexive, antisymmetric, transitive

\subseteq is a partial order



$\{0, 1, 2\} \subseteq \{2, 3, 4\}$? NO

$\{2, 3, 4\} \subseteq \{0, 1, 2\}$? NO

incomparable

Define \leq on \mathbb{R} by $a \leq b$ iff $a \subseteq b$

reflexive	$\forall a \in \mathbb{R} a \leq a$
transitive	$\forall a, b, c \in \mathbb{R} a \leq b \wedge b \leq c \rightarrow a \leq c$
antisymmetric	$\forall a, b \in \mathbb{R} a \leq b \wedge a \neq b \rightarrow \neg(b \leq a)$

DEF: A total order \leq on set A is a partial order on A such that for all $a, b \in A$, $a R b$ or $b R a$

\leq on \mathbb{R} is total order
b/c $\forall a, b \in \mathbb{R} a \leq b$ or $b \leq a$

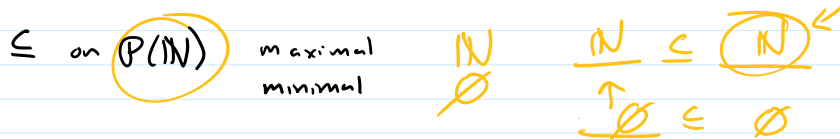
DEF: For a set A with partial order \leq ,

a is a maximal element means there is no $b \in A$ s.t. $a \leq b$ and $b \neq a$

a is the greatest element means for all $b \in A$ $b \leq a$ or $b = a$

a is a minimal element means there is no $b \in A$ s.t. $b \leq a$ and $b \neq a$

a is the least element means for all $b \in A$ $a \leq b$ or $b = a$



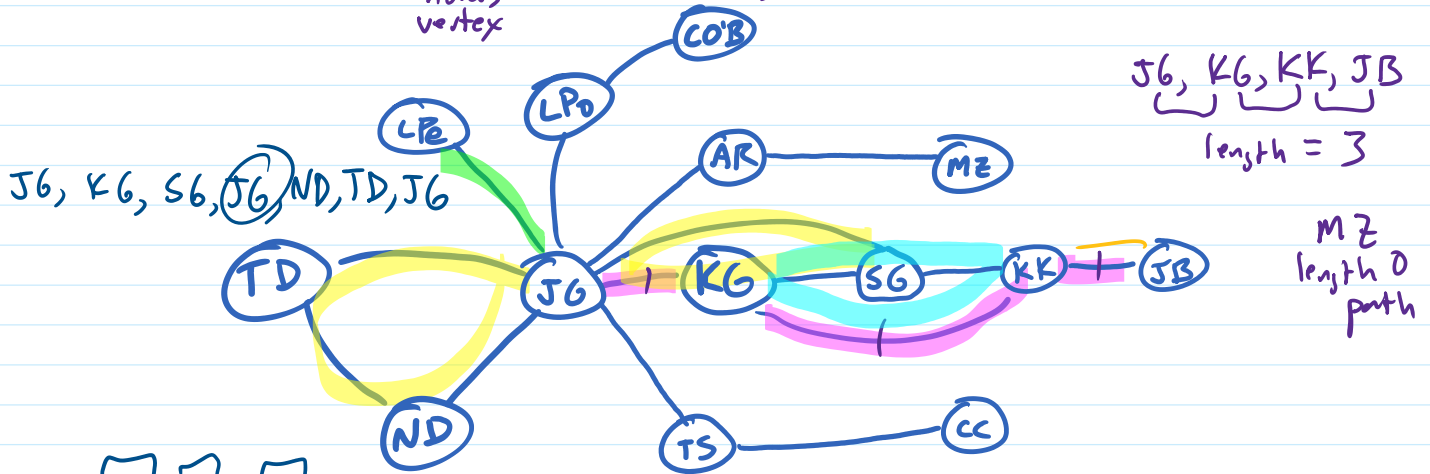
\leq on \mathbb{N}

maximal none!
minimal $0 \leq 0$

$$\underbrace{10^{100}}_{\leq} \leq \underbrace{10^{100} + 1}$$

Graphs

↳ representation of things and relationships between them
 people nodes vertex
 relationships edges



KG, SG, KK, KG
 SG, KK, KG, SG
 JG, KG, JG

simple graph: no self-loops
 no parallel edges

path: non-empty seq. of verts s.t.
 edge between adj verts in seq

simple path: path w/ no repeated verts
 of length ≥ 3

cycle: path beginning/ending at same place
 no repeated edges
 only repeated vertex is 1st/last

circuit: cycle w/ only 1st/last repeat relaxed

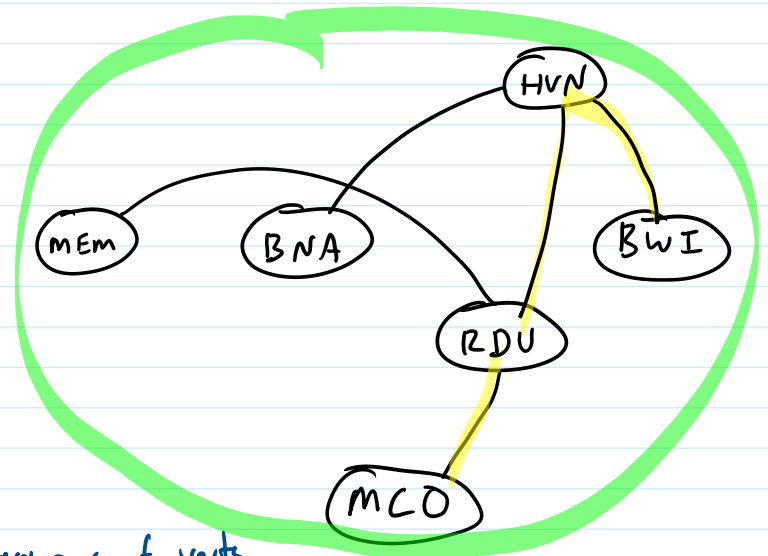
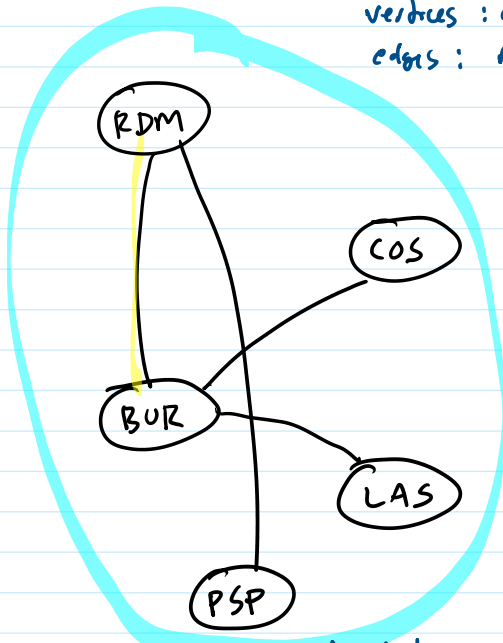
closed walk: path w/ same start/end

Airline Route Map

vertices : airports
edges : routes between airport

BWI

RDM

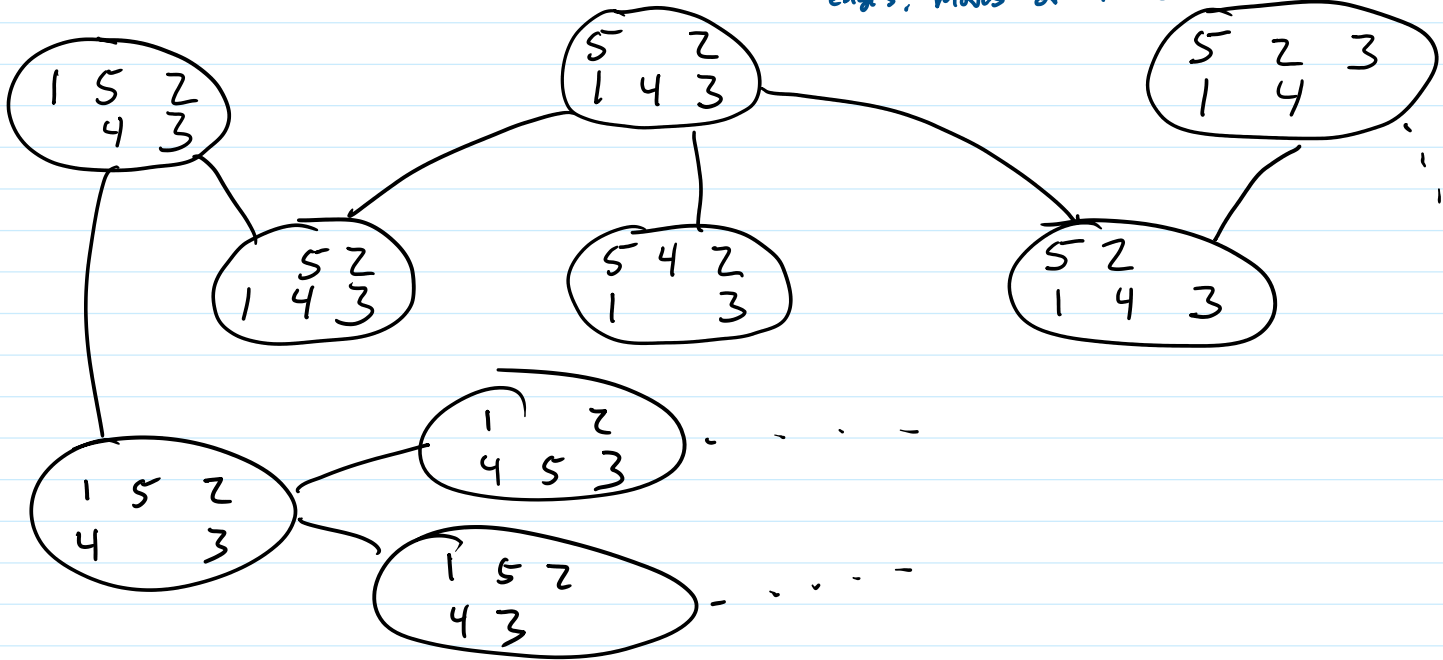


connected : path between every pair of vertices

connected component : maximal subset of vertices that is connected

5 puzzle

vertices: arrangements of tiles
edges: moves of 1 tile



Chomp



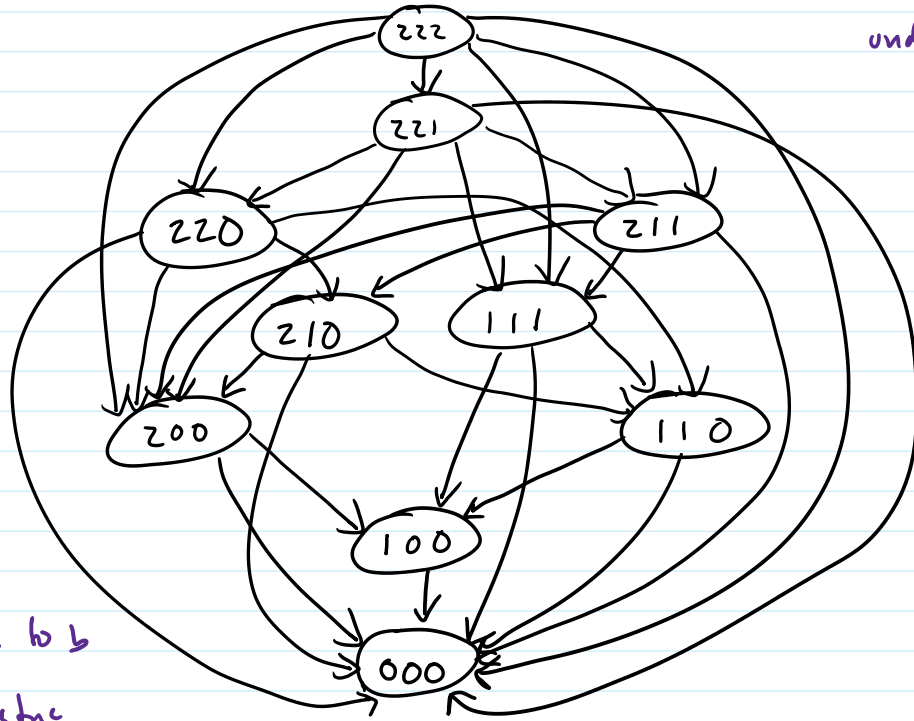
a

Edge $a \rightarrow b$

pick square in a
to leave state b

aRb iff
can move from a to b

chomp antisymmetric



undirected: irreflexive
symmetric

Paths

THM: If u, v connected in a graph $G=(V, E)$, then there is a path $u \rightsquigarrow v$ of length at most $n-1$.

$$n = n(V) \quad m = n(E)$$

↓ ↙

Proof: Let $G=(V, E)$ be a graph and suppose $u, v \in V$ are connected.

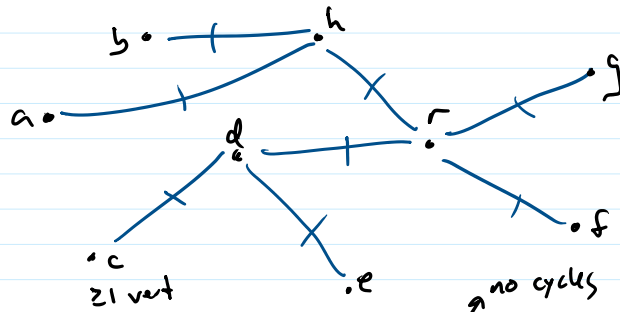
Then there is a shortest path v_0, \dots, v_k in G with $v_0 = u$ and $v_k = v$.

Two cases: 1) $k \leq n-1$

2) $k > n-1$

Tree

↳ connected, acyclic, undirected graph



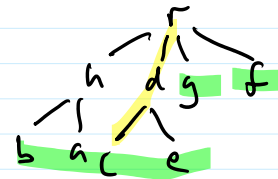
THM: Any non-empty, connected, acyclic, undirected graph has $n-1$ edges

THM: In a tree, there is a unique simple path between any pair of vertices

Rooted tree: tree w/ distinguished node (root)

parent of c = d

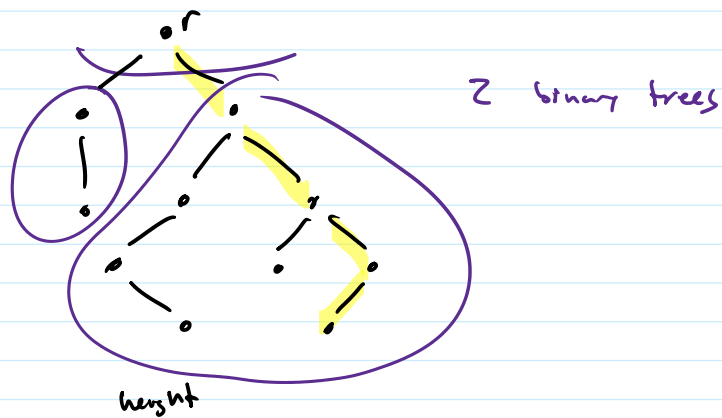
leaf: no children



height = 2

Binary Tree

Binary tree: a rooted tree in which every node has ≤ 2 children



THM: All binary trees have $\leq 2^{h+1} - 1$ nodes

Proof: We prove for all possible heights h , all binary trees of height h have $\leq 2^{h+1} - 1$ nodes

Base cases:

$(h=0)$ one node $2^{0+1} - 1 = 1$; $1 \leq 1$

Ind step: suppose $h > 0$ and for all i , $0 \leq i < h$, all trees of height i have $\leq 2^{i+1} - 1$ nodes

Let T be a tree of height h

$h_L, h_R \leq h-1$



$$\begin{pmatrix} n_L \\ n_R \end{pmatrix} \leq \begin{pmatrix} 2^{h_L+1} - 1 \\ 2^{h_R+1} - 1 \end{pmatrix} \leq \begin{pmatrix} 2^h - 1 \\ 2^h - 1 \end{pmatrix}$$

$$\begin{aligned} n &= n_L + n_R + 1 \leq 2^h - 1 + 2^h - 1 + 1 \\ &= \underline{2^{h+1} - 1} \end{aligned}$$