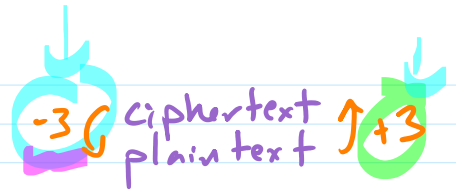


Cryptography

P D W K L V I X Q \$³⁶
M A T H I S F U N !³⁶



encryption key: used to transform plaintext → ciphertext
decryption key: used to transform ciphertext → plaintext

Public-Key Cryptography

Goal: allow anyone to send you an encrypted msg w/o prior key exchange
difficult for anyone other than recipient to decrypt message

public key: encryption key not secret - post anywhere

private key: decryption key secret; will be related to encryption key
but must be related in a way that makes it difficult to recover private from public

THM (Fermat's Little Theorem): For any prime p and $a \in \mathbb{Z}$ not a multiple of p ,
 then $a^{p-1} \equiv 1 \pmod{p}$

and $a^p \equiv a \pmod{p}$

$p=7$	$11 \cdot 13$ $p \cdot q$	n	1	2	3	4	5	6
		143	$143 \% 2 = 1$	$143 \% 5 = 3$	$143 \% 3 = 2$	$143 \% 7 = 3$	$143 \% 11 = 0$	
		$n^6 \pmod{7}$	1	$64 \equiv 1$	$3^6 = (3^2)^3 = 2^3 = 1$	$4^6 = 2^6 = 1$	$5^6 = 4^6 = 1$	$6^6 = 3^6 = 1$

common sizes of n : hundreds of digits

RSA

public key: $n = p \cdot q$ for primes p, q
 $e \in \mathbb{Z}$ s.t. $\gcd(e, \text{lcm}(p-1, q-1)) = 1$

private key: $d \in \mathbb{Z}$ s.t. $d \cdot e \equiv 1 \pmod{\text{lcm}(p-1, q-1)}$
 and $p > q$, d_p s.t. $d \equiv d_p \pmod{p-1}$, $0 \leq d_p < p-1$
 d_q s.t. $d \equiv d_q \pmod{q-1}$, $0 \leq d_q < q-1$

ciphertext

To encrypt: find c s.t. $m^e \equiv c \pmod{n}$

To decrypt: find m s.t. $c^d \equiv m \pmod{n}$

Example: choose $p = 37$ $q = 73$ so $n = 2701$

compute $\text{lcm}(36, 72) = 72$ (so a bad choice of p, q)
 choose $e = 5$ note $\gcd(5, 72) = 1$ better if $\gcd(p-1, q-1)$ is small

find d s.t. $de \equiv 1 \pmod{72}$ so $d = 29$ $145 = 5 \cdot 29$
 $5d \equiv 1 \pmod{72}$

$m = 1304$
 $M_D \rightarrow 1304 \rightarrow (1304)^5 = 1304 \cdot (1304)^4 \pmod{2701}$

$$= 1304 \cdot (1304^2)^2$$

$$\equiv 1304 \cdot (1487)^2$$

$$\equiv 1304 \cdot 1751$$

$$\equiv 959$$

$$29 = 16 + 8 + 4 + 1$$

$$x^{29} = x^{16} \cdot x^8 \cdot x^4 \cdot x$$

$$959^{29} = 959 \cdot (959^2)^2 \cdot ((959^2)^2)^2 \cdot (((959^2)^2)^2)^2$$

$$= aca \cdot 1741^2 \cdot (1741^2)^2 \cdot ((1741^2)^2)^2 \pmod{2701}$$

$$\begin{aligned}
&\equiv 959 \cdot 1341^2 \cdot (1341^2)^2 \cdot ((1341^2)^2)^2 \pmod{2701} \\
&\equiv 959 \cdot 2116 \cdot 2116^2 \cdot (2116^2)^2 \pmod{2701} \\
&\equiv 959 \cdot 2116 \cdot 1899 \cdot 1899^2 \pmod{2701} \\
&\equiv 959 \cdot 2116 \cdot 1899 \cdot 366 \equiv 1304 \pmod{2701}
\end{aligned}$$

MD

Decrypt

2066 →

using same private key

find plaintext that makes this work

$$2066^{29} \equiv \underline{\hspace{2cm}} \pmod{2701}$$

$$\equiv 2066 \cdot 2066^4 \cdot 2066^8 \cdot 2066^{16} \pmod{2701}$$

$$\equiv 2066 \cdot 2554 \cdot 1 \cdot 1$$

$$\equiv \frac{1511}{0K} \pmod{2701}$$

$\gcd(n_i, n_j) = 1$ for $i \neq j$

THM (Sun-tzu's Remainder Theorem): If n_1, \dots, n_k are pairwise co-prime integers, and $r_1, \dots, r_k \in \mathbb{Z}$ satisfy $0 \leq r_i < n_i$

then there is a unique integer r s.t.

$0 \leq r < n_1 \cdot \dots \cdot n_k$ and

$$\begin{cases} r \equiv r_1 \pmod{n_1} \\ r \equiv r_2 \pmod{n_2} \\ \vdots \\ r \equiv r_k \pmod{n_k} \end{cases}$$

(and r', r'' satisfy iff $r' \equiv r'' \pmod{n_1 \cdot \dots \cdot n_k}$)

There are
 $m \equiv r_p \pmod{p}$
 $m \equiv r_g \pmod{g}$
 if $m^{de} \equiv r_p \pmod{p}$ and $m^{de} \equiv r_g \pmod{g}$ (\rightarrow dir of n iff)
 then $m^{de} \equiv m \pmod{p \cdot g}$
 $\gcd(3, 5) = 1$

$\gcd(3, 4) = 1$ $\gcd(4, 5) = 1$

Ex: $n_1 = 3$ $n_2 = 4$ $n_3 = 5$

	0	1	2	3	4	5	$3 \cdot 4 \cdot 5 - 1$
$n \pmod 3$	0	1	2	0	1	2
$n \pmod 4$	0	0	1	2	3	0
$n \pmod 5$	0	1	2	3	4	0

THM: Suppose p, g are primes, $n = p \cdot g$,

d, e are integers s.t. $\gcd(e, \text{lcm}(p-1, g-1)) = 1$, $d \cdot e \equiv 1 \pmod{\lambda(n)}$

$(m^e)^d = m^{de}$ Then $m^{de} \equiv m \pmod{n}$

Proof: Suppose p, g, n, d, e are as given

It suffices to show $m^{de} \equiv m \pmod{p}$ and $m^{de} \equiv m \pmod{g}$ (by Sun-tzu)

$m^{de} \equiv m \pmod{p}$: note $p-1 \mid \text{lcm}(p-1, g-1)$

Two cases:

1) $m \equiv 0 \pmod{p}$
 then $m^{de} = 0^{de} = 0 = m \pmod{p}$

also

$d \cdot e \equiv 1 \pmod{\lambda(n)}$
 so $d \cdot e - 1 = \lambda(n) \cdot l$ for some $l \in \mathbb{Z}$

$d \cdot e - 1 = (p-1) \cdot k \cdot l$

$d \cdot e = 1 + (p-1) \cdot h$

$m^{de} = m^{1 + (p-1) \cdot h} = m \cdot (m^{p-1})^h$

$$m^{de} = m^{1+(p-1)h} = m \cdot (m^{p-1})^h$$

Fermat's Little Thm \longrightarrow $\equiv m \cdot 1^h \pmod{p}$
Says $m^{p-1} \equiv 1 \pmod{p}$
since we are in the $m \not\equiv 0 \pmod{p}$ case $\equiv m \pmod{p}$

$m^{de} \equiv m \pmod{p}$: similar