

# Counting with Population Protocols

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Yale University



*Inria*

# Population protocol model

Introduced by **Angluin, Aspnes, Diamadi, Fischer and Peralta in 2004**<sup>1</sup>

A model of distributed systems with minimal set of assumptions

- 1 Finite-state agents / automata
  - 2 Agents have no identities
  - 3 Communication is occasionally possible between agents
- **Uniformity** All the agents execute the same code and this code is independent of the population size
- **Anonymity**: No room for agents to store a unique identifier

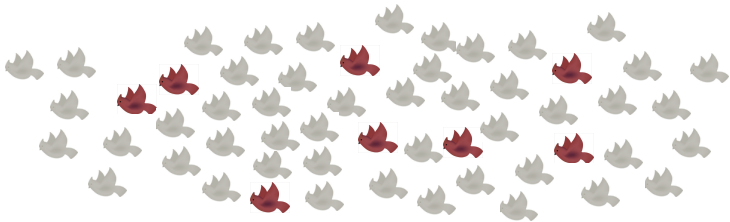
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<sup>1</sup>D. Angluin, J. Aspnes, Z. Diamadi, M. J. Fischer, R. Peralta, "Computation in networks of passively mobile finite-state sensors", PODC 2004: 290-299

# A motivating example: birds [Angluin04]

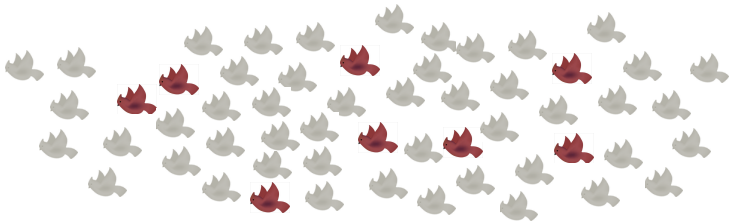
A flock of birds on which are strapped identical sensors  
When birds are close together, their sensors can interact and compute simple functions

- Decide if at least five birds have elevated body temperature
- Decide if at least five percent of the birds have elevated body temperature
- Compute the proportion of birds that have elevated temperature
- ...



# Agenda

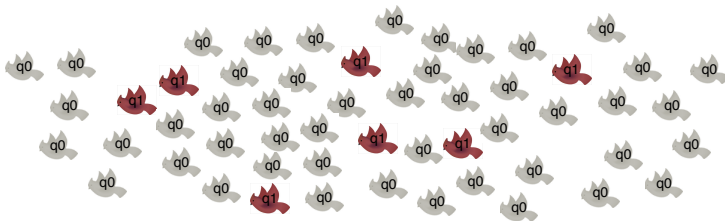
- Motivation to study the population protocol model
- Model
- The counting problem
- Performance evaluation
- Conclusion and future works





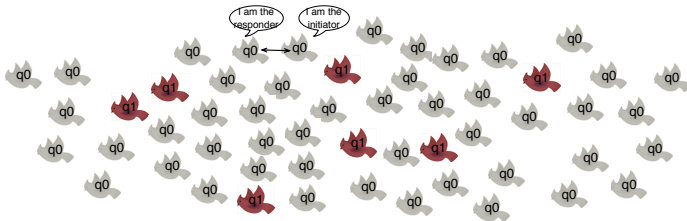
# The population protocol model (informal)

- Collection of **finitely many** finite-state agents
- Agents are **indistinguishable** (identical program + no identity)
- Agents enter their **initial state** by applying an input function  $I$  to a **finite set**  $X$  of input values



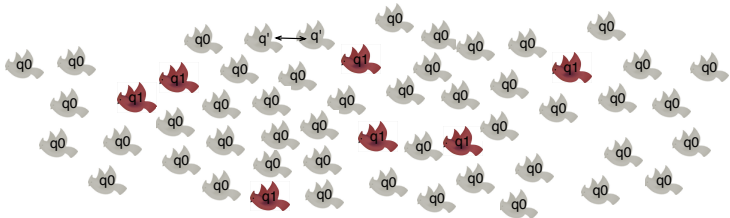
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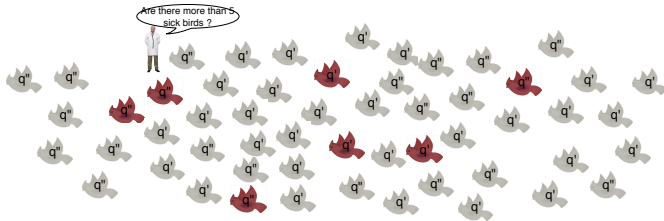
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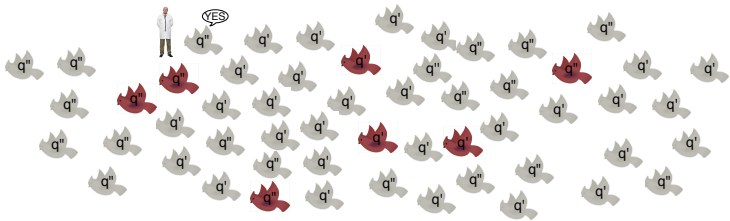
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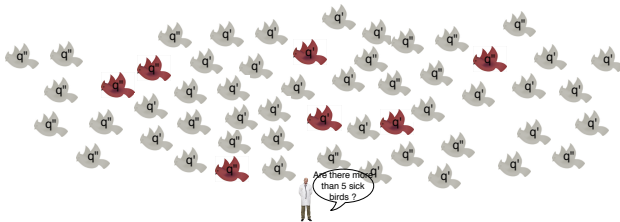
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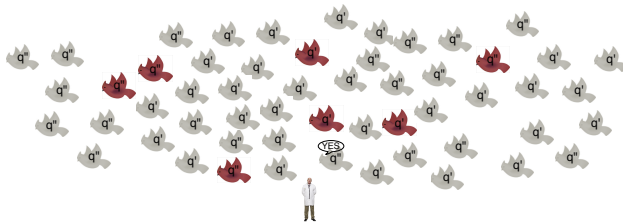
All agents eventually **converge to a correct common or distributed output value**



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# The population protocol model (informal)

- A population **configuration**  $C$  specifies the state of each agent
- **Fairness condition**: enforces that any reachable configuration is eventually reached.
- A stable computation is a infinite fair sequence of configurations that **converges to a correct common or distributed output value**
  - Convergence is a global property: agents generally do not know that convergence has been reached.
  - With suitable stochastic assumptions, it is possible to determine the number of interactions until the output stabilizes



# A motivating example: birds [Angluin 04]

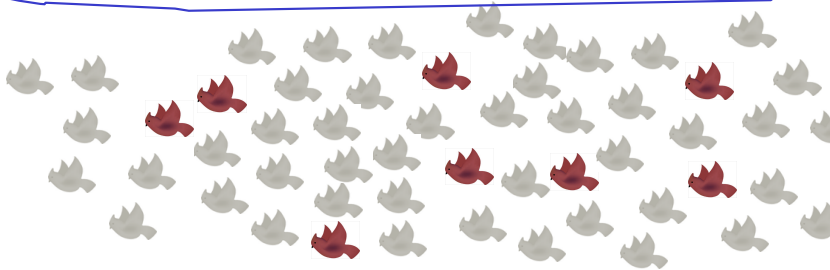
A flock of birds on which are strapped identical sensors

When birds are close together, their sensors can interact and compute simple functions

Decide if **at least five birds** have elevated body temperature ?

Decide if **at least five percent of the birds** have elevated body temperature?

Compute the **proportion of birds** that have elevated temperature ?



# Computational power of population protocols

- A predicate can be seen as a function that returns `true` or `false`
- Predicates can be written as  $P(x_1, x_2, \dots, x_k)$ , where
  - $k$  = number of possible initial states
  - $x_i$  = number of agents starting in the  $i$ -th state
- Examples:
  - the "count-to-5" bird protocol:  $P(x_1, x_2) = \text{true}$  iff  $x_1 \geq 5$
  - the "majority" bird protocol:  $P(x_1, x_2) = \text{true}$  iff  $x_1 \geq x_2$
  - the "parity" bird protocol:  $P(x_1, x_2) = \text{true}$  iff  $x_1 \equiv 0 \pmod{2}$

## Theorem (Angluin et al. 2007)

*A predicate is computable by a population protocol if and only if it is in one of the following forms<sup>a</sup>:*

- $\sum_{i=1}^k c_i x_i \geq a$
- $\sum_{i=1}^k c_i x_i \equiv a \pmod{b}$ ,  
where  $a_i$ ,  $b$  and  $c_i$ 's are integer constant
- *And every boolean combination of these predicates*

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<sup>a</sup>D. Angluin, J. Aspnes, D. Eisenstat, E. Ruppert, " The computational power of population protocols", Distributed Computing 20(4): 279-304 (2007)

# Our contribution: counting the exact percentage of sick and safe birds

## Problem

*Design of a population protocol that exactly determines the difference between the percentage of sick birds and the percentage of safe birds.*

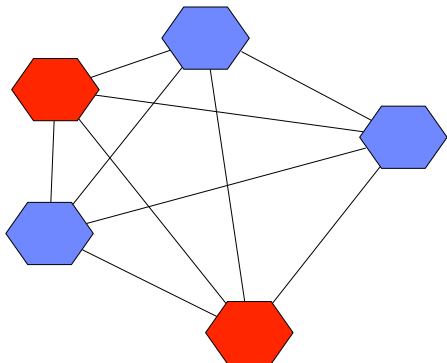
# Our algorithm to compute the exact difference between the percentage of #red and #blue

## Algorithm

- *Finite input alphabet*  $X = \{\text{sick}, \text{safe}\}$
- *Input function*  $I: I(\text{sick}) = m, I(\text{safe}) = -m$
- *Finite set of states*  $Q = \{-m, -m + 1, \dots, m - 1, m\}$
- *Output function*  $O: O(q) = \lfloor 100q/m + 1/2 \rfloor$
- *Transition function*  
 $f: (q_1, q_2) \rightarrow (\lfloor (q_1 + q_2)/2 \rfloor, \lceil (q_1 + q_2)/2 \rceil)$

## Problem

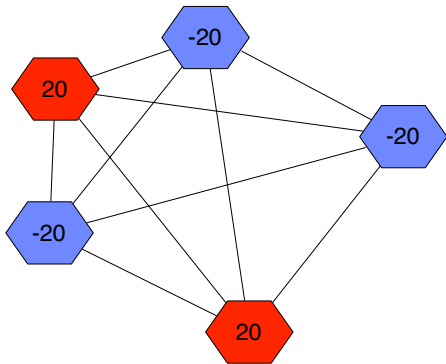
*What is the difference between the percentage of sick and the percentage of safe birds?*



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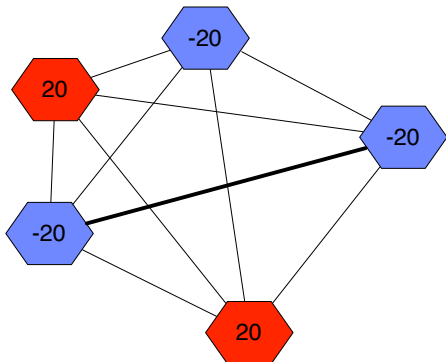
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- $\sum_{i=1}^5 = \text{constant} = -20$ .  
As you will see this is an invariant of the protocol



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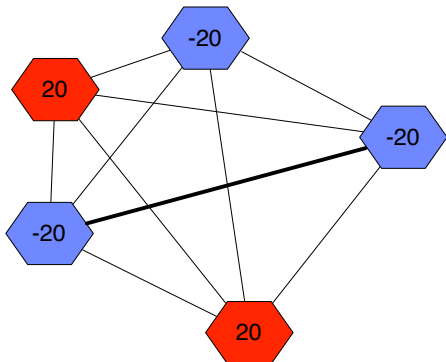
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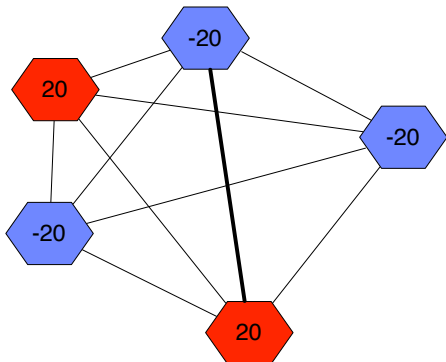
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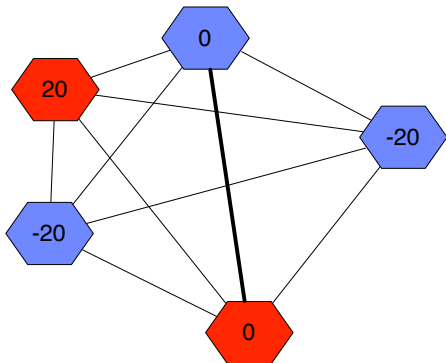
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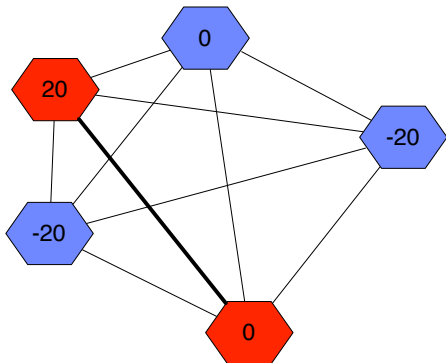
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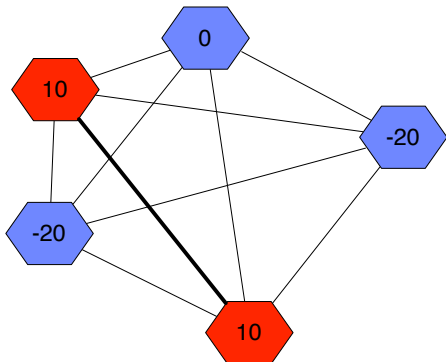
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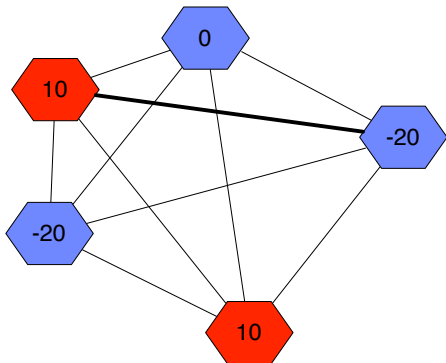
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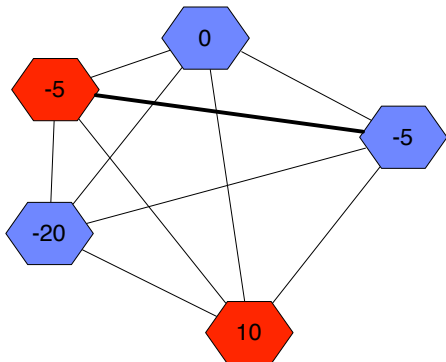
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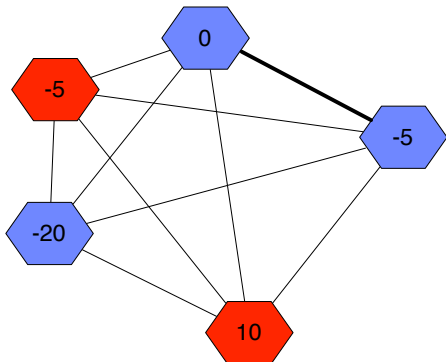
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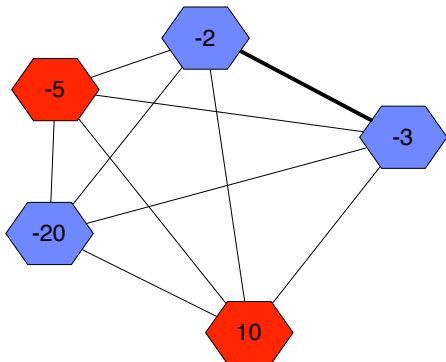
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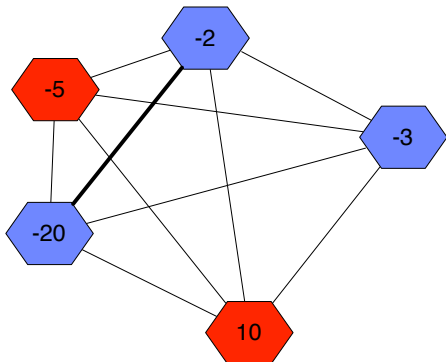
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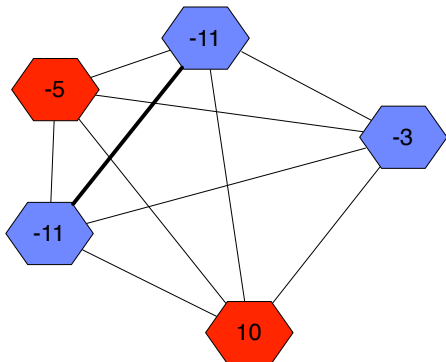
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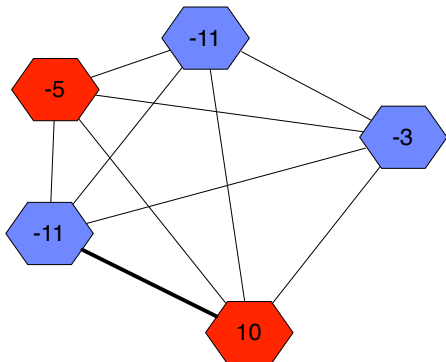
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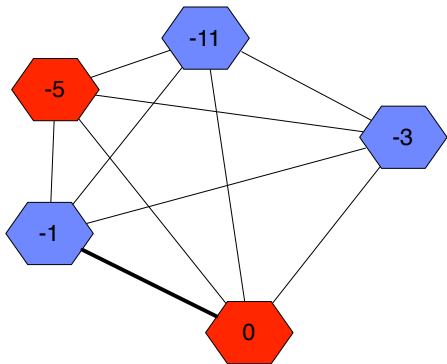
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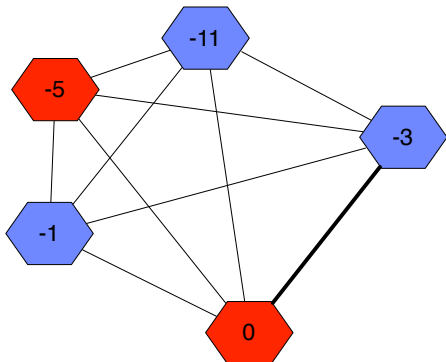
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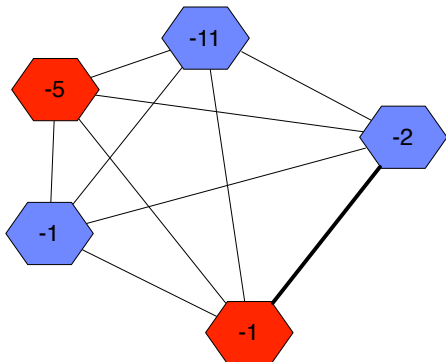
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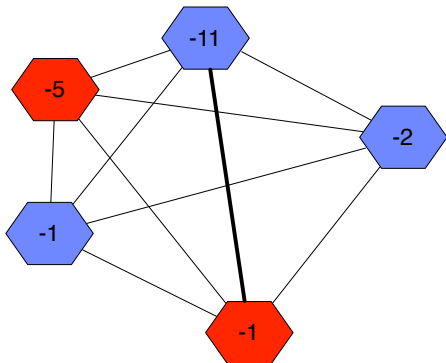
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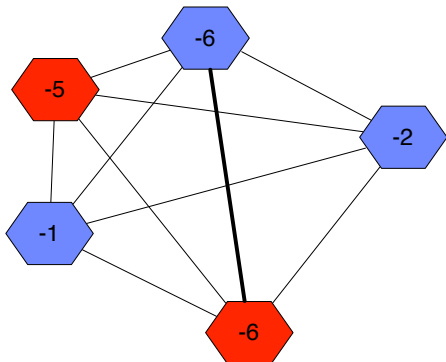
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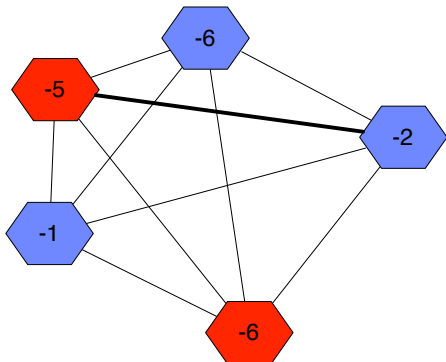
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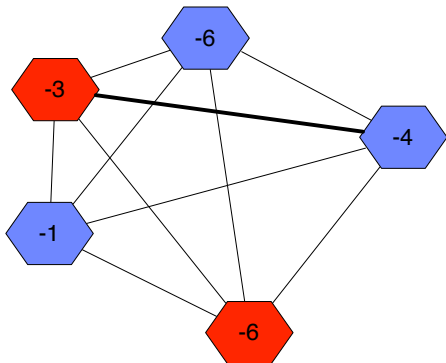
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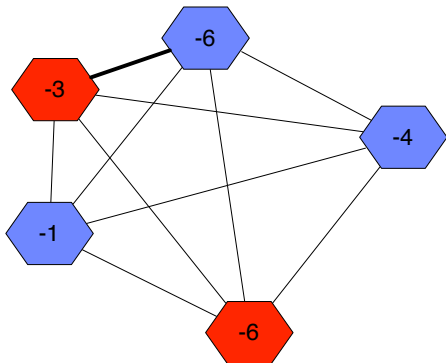
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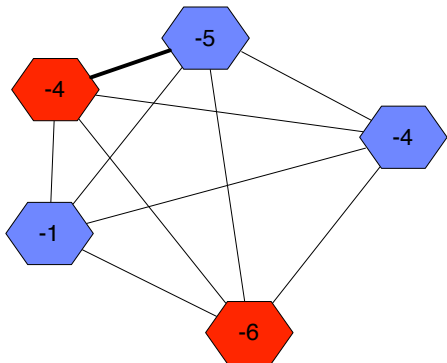
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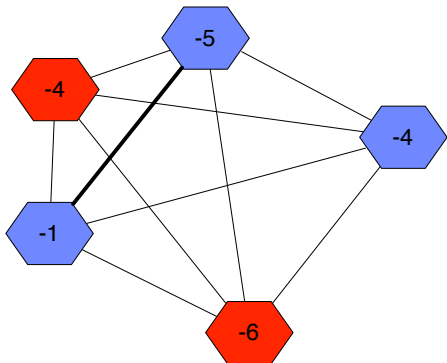
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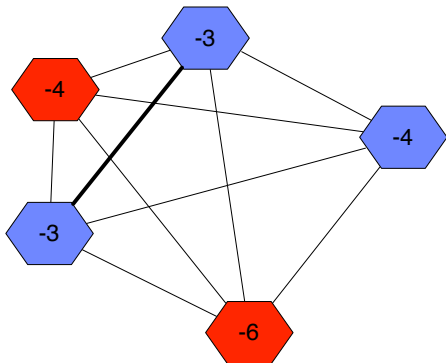
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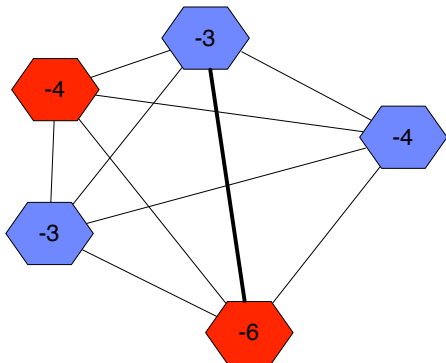
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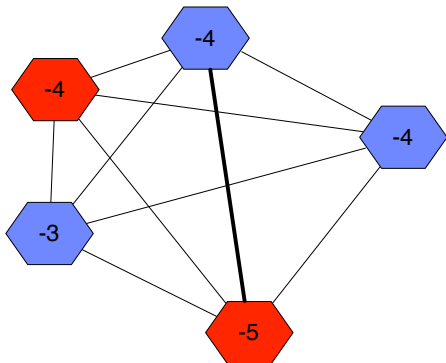
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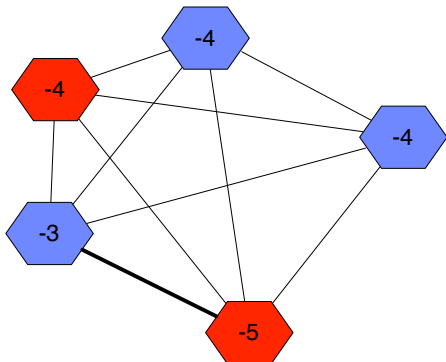
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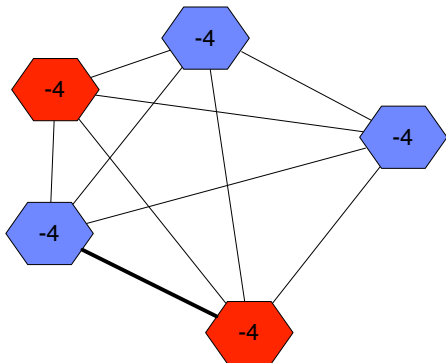
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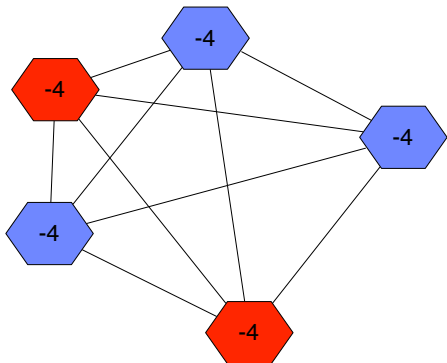
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$$\lfloor 100q/m + 1/2 \rfloor = \lfloor 100 \times (-4)/20 + 1/2 \rfloor = -20$$



# Computing with randomized interactions

- The choice of the interacting agents is a **non-deterministic** choice
- **Probabilistic assumptions**: interactions are orchestrated by a uniform fair random scheduler:
  - At each discrete time  $t$ , any two agents  $i$  and  $j$  are randomly chosen from a uniform distribution  $p_{i,j}(t)$ , with

$$p_{i,j}(t) = \frac{1}{n(n-1)}$$

# How long does it take for each agent to converge ?

## Main ideas of the convergence proof.

- We show that at each step, the sum of the agents' state is constant. For every  $t \geq 0$ ,

$$\sum_{i=1}^n C_t^{(i)} = \sum_{i=1}^n C_0^{(i)}$$

- Let  $\ell = \frac{1}{n} \sum_{i=1}^n C_t^{(i)}$  and  $L = (\ell, \dots, \ell)$ , we show that

$$E \left( \|C_t - L\|^2 \right) = \left( 1 - \frac{1}{n-1} \right)^t E \left( \|C_0 - L\|^2 \right) + \frac{n}{4}.$$



# How long does it take for each agent to converge ?

## Main ideas of the convergence proof (continued).

- Let  $\kappa = (\#red - \#blue)/100$
- For all  $\delta \in (0, 1)$ ,  $m = \lceil \sqrt{2}n^{3/2}/\sqrt{\delta} \rceil$  and for all  $t \geq (n-1) \left( 5 \ln 2 + 3 \ln n - \ln \delta + \frac{2}{m-1} \right)$ , we show that

$$P\{O(C_t^{(i)}) = \kappa, \text{ for all } i = 1, \dots, n\} \geq 1 - \delta$$

- Thus the parallel convergence time to get  $\kappa$  with any high probability is  $O(\log n)$
- The knowledge of the population size  $n$  allows us to directly derive  $\#red - \#blue$



# Performance evaluation

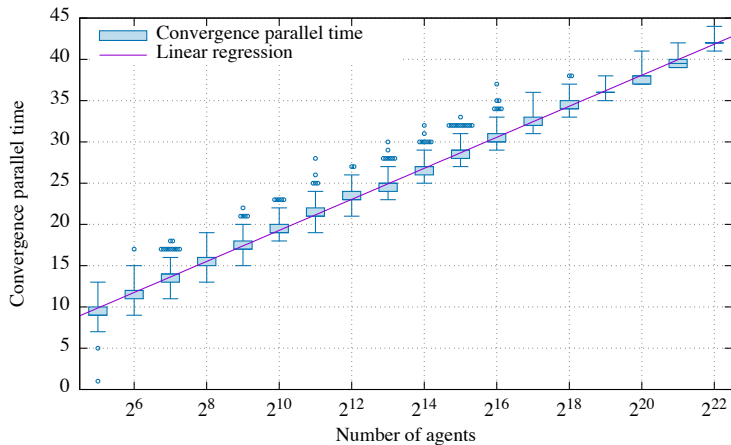


Figure : Evolution of the configuration vector for a conserved advantage  
#red - #blue equal to  $3n/5$ . Settings:  $n = 2^{22} = 4,19 \times 10^6$ .



# Performance evaluation

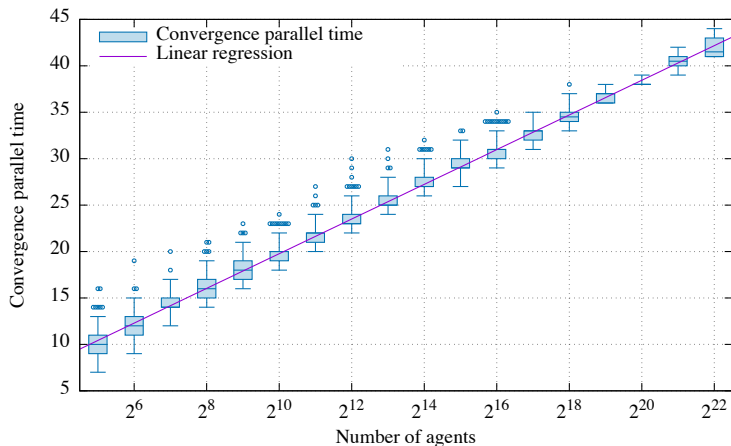


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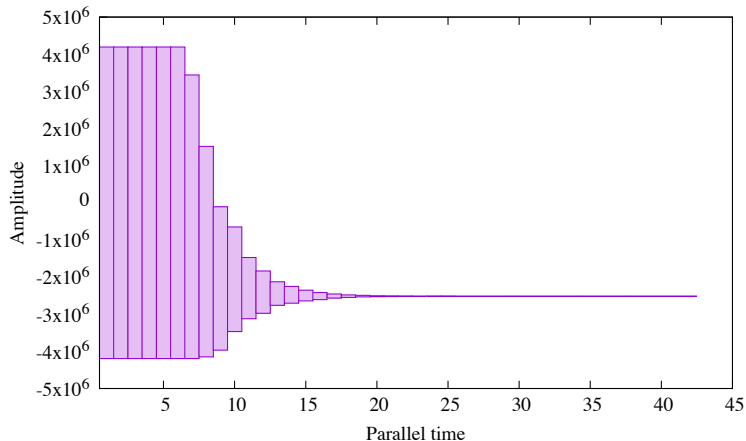


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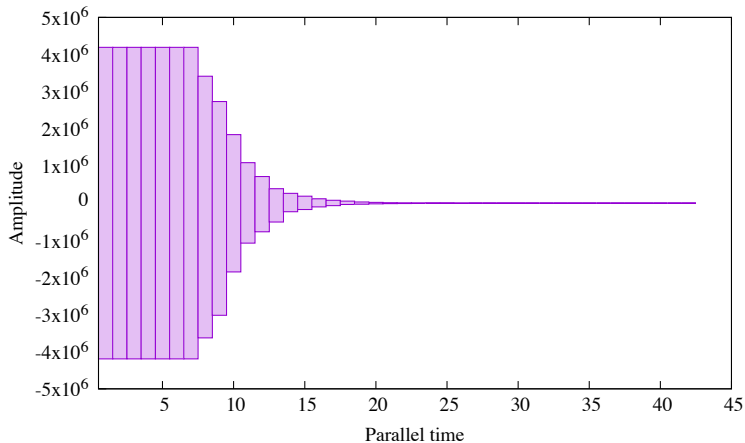


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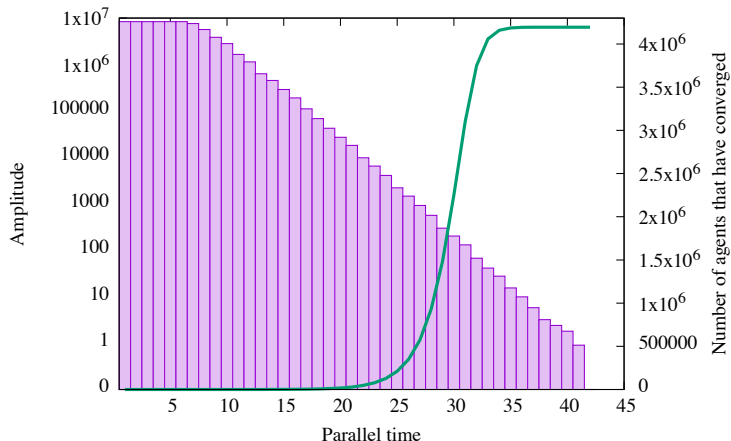


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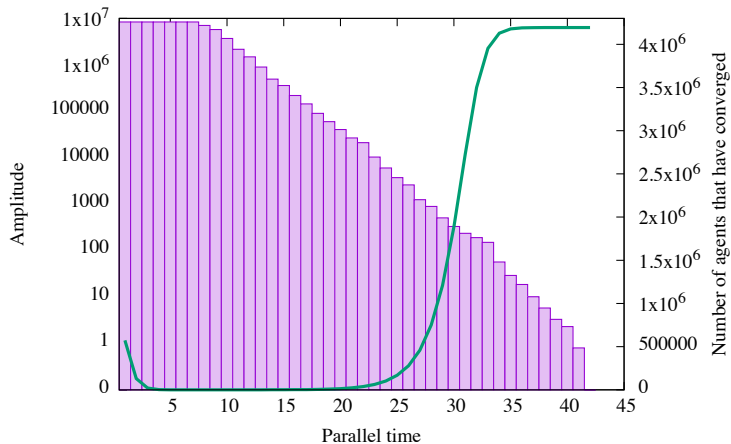


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# Related works: Computation of the majority, i.e., whether $\#red > \#blue$

- [DV12]<sup>2</sup>, [MNRS14]<sup>3</sup>:
  - Four-state protocol with an expected convergence parallel time in  $O(\log n)$
  - When  $\#red \simeq \#blue$ , the convergence parallel time is infinite

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<sup>2</sup>[DV12] M. Draief and M. Vojnovic, “Convergence speed of binary interval consensus”, SIAM Journal on Control and Optimization, 50(3):1087:1097, 2012

<sup>3</sup>[MNRS14] G. Mertzios, et al., “Determining majority in networks with local interactions and very small local memory”, ICALP, pp 871-882, 2014

# Related works: Computation of the majority, i.e., whether $\#red > \#blue$

- [AAE07]<sup>4</sup>, [PVV09]<sup>5</sup>:
  - Three-state protocol with an expected convergence parallel time in  $O(\log n)$
  - Only when  $\#red - \#blue$  is in  $O(\sqrt{\log n})$

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<sup>4</sup>[AAE07] D. Angluin, J. Aspnes and D. Eisenstat, “A simple population protocol for fast Robust approximation majority”, Distributed Computing, 20(4):279-305, 2007

<sup>5</sup>[PVV09] E. Perron, D. Vasudevan and M. Vojnovic, “Using three states for binary consensus on complete graphs”, INFOCOM, pp 2527-3435, 2009

# Related works: Computation of the majority, i.e., whether $\#red > \#blue$

- [AGV15]<sup>6</sup>:
  - $(\log n)$ -state protocol with an expected convergence parallel time in  $O(\log n)$
  - Whatever the difference between  $\#red$  and  $\#blue$
  - The authors show that a convergence in  $O(\log n)$  interactions in expectation is a lower bound.

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<sup>6</sup>[AFV15] D. Alistarh, R. Gelashvili, and M. Vojnovic, “Fast and exact majority in population protocols”, Technical report MSR-TR-2015-13, Microsoft research, 2015



- The population protocol model: A simple but powerful enough model to describe distributed systems
- The exact computation power of these protocols has been determined
- We have proposed simple proofs to show that the convergence time for the counting problem is logarithmic in the size of the population